

CS 58500 – Theoretical Computer Science Toolkit

Lecture 3 (01/27)

Concentration Inequality II

https://ruizhezhang.com/course_spring_2026.html



Today's Lecture

- Chernoff Bound (Additive Form)
- Chernoff Bound (Multiplicative Form)
- Application: Discrepancy Theory

Remark: in many TCS papers,

“Chernoff bound” \approx exponential convergence \approx Chernoff/Hoeffding/Bernstein inequality

Question

Suppose you have a coin with **unknown biased** probability $p \in [0, 1]$. How many flips to estimate p with approximation error $\pm\epsilon$?

- Flip n times and output $\hat{p} := \# \text{ Heads}/n$
- Suppose we want to achieve the guarantee

$$\Pr[|\hat{p} - p| \leq \epsilon] \geq 1 - \delta$$

- How small can we choose n ?



Question

Suppose you have a coin with **unknown biased** probability $p \in [0, 1]$. How many flips to estimate p with approximation error $\pm\epsilon$?



- Define **iid** random variables X_1, X_2, \dots, X_n such that

$$X_i := \begin{cases} 1 & \text{with probability } p \\ 0 & \text{with probability } 1 - p \end{cases}$$

- For $Z := X_1 + \dots + X_n$, how large should n be such that

$$\Pr[|Z - pn| \geq \epsilon n] \leq \delta$$

- $\mathbb{E}[Z] = pn$
- $\text{Var}[X_i] = p - p^2$ and $\text{Var}[Z] = np(1 - p)$

Question

Suppose you have a coin with **unknown biased** probability $p \in [0, 1]$. How many flips to estimate p with approximation error $\pm\epsilon$?

- $\mathbb{E}[Z] = pn$ and $\text{Var}[Z] = np(1 - p)$
- Chebyshev inequality:

$$\Pr[|Z - \mathbb{E}[Z]| \geq t] \leq \frac{\text{Var}[Z]}{t^2}$$

$$\Pr[|Z - pn| \geq \epsilon n] \leq \frac{np(1 - p)}{\epsilon^2 n^2} = \frac{p(1 - p)}{\epsilon^2 n} \leq \delta$$

$$n \geq \frac{p(1 - p)}{\epsilon^2} \frac{1}{\delta}$$



Chernoff Bound

- Let $X'_i := X_i - \mathbb{E}[X_i]$ be the centered version of X_i and $Z' := X'_1 + \dots + X'_n = Z - \mathbb{E}[Z]$
- Consider the **moment generating function**:

$$\mathbb{E}[e^{\theta Z'}] = \mathbb{E}[e^{\theta \sum_i X'_i}] = \mathbb{E}\left[\prod_i e^{\theta X'_i}\right] = \prod_i \mathbb{E}[e^{\theta X'_i}]$$

- By Markov inequality,

$$\Pr[Z' \geq \epsilon n] = \Pr[e^{\theta Z'} \geq e^{\theta \epsilon n}] \leq \frac{\mathbb{E}[e^{\theta Z'}]}{e^{\theta \epsilon n}} = \frac{\prod_i \mathbb{E}[e^{\theta X'_i}]}{e^{\theta \epsilon n}}$$

- This inequality holds for any $\theta > 0$. So, we can optimize over t to get the tightest bound:

$$\Pr[Z' \geq \epsilon n] \leq \inf_{\theta > 0} e^{-\theta \epsilon n} \prod_i \mathbb{E}[e^{\theta X'_i}]$$

Chernoff Bound

Lemma (MGF bound). If $a \leq X \leq b$ then for every $\theta \in \mathbb{R}$ we have

$$\mathbb{E}[e^{\theta(X - \mathbb{E}[X])}] \leq e^{\theta^2(b-a)^2/8}$$

Proof.

- Wlog, we may assume $\mathbb{E}[X] = 0$ and $\theta \geq 0$
- Define the log-MGF:

$$\psi(\theta) := \log \mathbb{E}[e^{\theta X}]$$

- $\psi(0) = \log \mathbb{E}[1] = 0$
- We compute the derivatives:

$$\psi'(\theta) = \frac{\mathbb{E}[Xe^{\theta X}]}{\mathbb{E}[e^{\theta X}]}, \quad \psi''(\theta) = \frac{\mathbb{E}[X^2 e^{\theta X}]}{\mathbb{E}[e^{\theta X}]} - \left(\frac{\mathbb{E}[Xe^{\theta X}]}{\mathbb{E}[e^{\theta X}]} \right)^2$$

Chernoff Bound

Lemma (MGF bound). If $a \leq X \leq b$ then for every $\theta \in \mathbb{R}$ we have

$$\mathbb{E}[e^{\theta(X - \mathbb{E}[X])}] \leq e^{\theta^2(b-a)^2/8}$$

Proof.

- We compute the derivatives:

$$\psi'(\theta) = \frac{\mathbb{E}[X e^{\theta X}]}{\mathbb{E}[e^{\theta X}]}, \quad \psi''(\theta) = \frac{\mathbb{E}[X^2 e^{\theta X}]}{\mathbb{E}[e^{\theta X}]} - \left(\frac{\mathbb{E}[X e^{\theta X}]}{\mathbb{E}[e^{\theta X}]} \right)^2 \leftarrow \text{Looks like a variance}$$

- $\psi'(0) = \mathbb{E}[X] = 0$ by assumption
- Let $\mu := \text{Law}(X)$. Then we can define its **exponential tilt** μ_θ :

$$\frac{\mu_\theta(dx)}{\mu(dx)} := \frac{e^{\theta x}}{\mathbb{E}[e^{\theta X}]}$$

- You can easily check that μ_θ is indeed a probability distribution for any $\theta \geq 0$

Chernoff Bound

Lemma (MGF bound). If $a \leq X \leq b$ then for every $\theta \in \mathbb{R}$ we have

$$\mathbb{E}[e^{\theta(X - \mathbb{E}[X])}] \leq e^{\theta^2(b-a)^2/8}$$

Proof.

- We compute the derivatives:

$$\psi'(\theta) = \mathbb{E}_{\mu_\theta}[X], \quad \psi''(\theta) = \text{Var}_{\mu_\theta}[X]$$

- Recall the **range bound for variance**:

$$\psi''(\theta) = \text{Var}_{\mu_\theta}[X] \leq (b - a)^2/4$$

- Now, we integrate ψ'' twice from 0 to θ and using the fact that $\psi(0) = \psi'(0) = 0$:

$$\int_0^\theta \psi''(\lambda) d\lambda = \psi'(\theta), \quad \int_0^\theta \psi'(\lambda) d\lambda = \psi(\theta)$$

$$\Rightarrow \psi(\theta) = \int_0^\theta \int_0^s \psi''(\lambda) d\lambda ds \leq \int_0^\theta \int_0^s \frac{1}{4} (b - a)^2 d\lambda ds = \int_0^\theta \frac{s}{4} (b - a)^2 ds = \frac{\theta^2}{8} (b - a)^2$$



Chernoff Bound

Lemma (MGF bound). If $a \leq X \leq b$ then for every $\theta \in \mathbb{R}$ we have

$$\mathbb{E}[e^{\theta(X - \mathbb{E}[X])}] \leq e^{\theta^2(b-a)^2/8}$$

- This lemma implies that

$$\begin{aligned} \Pr[Z' \geq \epsilon n] &\leq \inf_{\theta > 0} e^{-\theta \epsilon n} \prod_i \mathbb{E}[e^{\theta X'_i}] \leq \inf_{\theta > 0} e^{-\theta \epsilon n} (e^{\theta^2(b-a)^2/8})^n \\ &= \exp\left(\inf_{\theta > 0} -\theta \epsilon n + \theta^2(b-a)^2 n/8\right) \end{aligned}$$

- The quadratic function is minimized at $\theta = \frac{4\epsilon}{(b-a)^2}$:

$$\Pr[Z - \mathbb{E}[Z] \geq \epsilon n] = \Pr[Z' \geq \epsilon n] \leq e^{-\frac{2\epsilon^2}{(b-a)^2}n}$$

- What about the other direction (the lower tail)?

Chernoff Bound

Lemma (MGF bound). If $a \leq X \leq b$ then for every $\theta \in \mathbb{R}$ we have

$$\mathbb{E}[e^{\theta(X - \mathbb{E}[X])}] \leq e^{\theta^2(b-a)^2/8}$$

- For the lower tail:

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \leq -\epsilon n] &= \Pr[Z' \leq -\epsilon n] \\ &= \Pr[e^{\theta Z'} \geq e^{-\theta \epsilon n}] \quad (\theta \leq 0) \\ &\leq \inf_{\theta \leq 0} e^{+\theta \epsilon n} \prod_i \mathbb{E}[e^{\theta X'_i}] \\ &= \exp\left(\inf_{\theta \leq 0} \theta \epsilon n + \theta^2(b-a)^2 n/8\right) \\ &= \exp\left(-\frac{2\epsilon^2}{(b-a)^2} n\right)\end{aligned}$$

Chernoff Bound (Additive Form)

Theorem (Hoeffding's inequality).

Let X_1, \dots, X_n be **independent** real random variables such that $a_i \leq X_i \leq b_i$ for all $i \in [n]$.

Let $Z := \sum_{i=1}^n X_i$ and introduce the variance proxy

$$v := \frac{1}{4} \sum_{i=1}^n (b_i - a_i)^2.$$

Then

$$\Pr[|Z - \mathbb{E}[Z]| \geq t\sqrt{v}] \leq 2e^{-t^2/2} \quad \forall t \geq 0$$

➤ For our biased coin problem, $X_i \in \{0,1\}$ and thus, $v = n/4$

$$\Pr[|Z - \mathbb{E}[Z]| \geq \epsilon n] \leq 2e^{-2\epsilon^2 n} \leq \delta \quad \Rightarrow \quad n \geq \frac{1}{2\epsilon^2} \log\left(\frac{2}{\delta}\right)$$

Chebyshev:

$$n \geq \frac{p(1-p)}{\epsilon^2} \frac{1}{\delta}$$

Improvement?

In the whole proof, there are only two inequalities:

1. By Markov inequality, we obtain the so-called **Laplace transform**:

$$\Pr[Z - \mathbb{E}[Z] \geq t] \leq \inf_{\theta \geq 0} e^{-t\theta} \prod_i \mathbb{E}[e^{\theta(X_i - \mathbb{E}[X_i])}],$$

$$\Pr[Z - \mathbb{E}[Z] \leq -t] \leq \inf_{\theta \leq 0} e^{t\theta} \prod_i \mathbb{E}[e^{\theta(X_i - \mathbb{E}[X_i])}]$$

Cramér's theorem:

This is tight for iid sum

2. **MGF bound**: If $a \leq X \leq b$ then for every $\theta \in \mathbb{R}$ we have

$$\mathbb{E}[e^{\theta(X - \mathbb{E}[X])}] \leq e^{\theta^2(b-a)^2/8}$$

- In our case, X_1, \dots, X_n are iid Bernoulli random variables

$$\mathbb{E}[e^{\theta(X_i - p)}] = pe^{\theta(1-p)} + (1-p)e^{-\theta p} = e^{-\theta p}(1 + (e^\theta - 1)p)$$

Improvement?

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \geq t] &\leq \inf_{\theta \geq 0} e^{-t\theta} \prod_i \mathbb{E}[e^{\theta(X_i - \mathbb{E}[X_i])}] \\ &= \inf_{\theta \geq 0} e^{-t\theta} e^{-np\theta} (1 + (e^\theta - 1)p)^n \\ &\leq \inf_{\theta \geq 0} e^{-t\theta} e^{-np\theta} e^{np(e^\theta - 1)} \\ &= \exp\left(-np + \inf_{\theta \geq 0} -(t + np)\theta + npe^\theta\right)\end{aligned}$$

- The minimizer is at $\theta = \log\left(\frac{t+np}{np}\right)$:

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \geq t] &\leq \exp\left(-np - (t + np) \log\left(\frac{t + np}{np}\right) + (t + np)\right) \\ &= \exp\left(-(t + np) \log\left(\frac{t + np}{np}\right) + t\right)\end{aligned}$$

Improvement?

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \geq t] &\leq \inf_{\theta \geq 0} e^{-t\theta} \prod_i \mathbb{E}[e^{\theta(X_i - \mathbb{E}[X_i])}] \\ &= \inf_{\theta \geq 0} e^{-t\theta} e^{-np\theta} (1 + (e^\theta - 1)p)^n \\ &\leq \inf_{\theta \geq 0} e^{-t\theta} e^{-np\theta} e^{np(e^\theta - 1)} \\ &= \exp\left(-np + \inf_{\theta \geq 0} -(t + np)\theta + npe^\theta\right)\end{aligned}$$

- The minimizer is at $\theta = \log\left(\frac{t+np}{np}\right)$:

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \geq tnp] &\leq \exp\left(-np - (tnp + np) \log\left(\frac{tnp + np}{np}\right) + (tnp + np)\right) \\ &= \exp(-(1+t)np \log(1+t) + tnp) \\ &= \left(\frac{e^t}{(1+t)^{1+t}}\right)^{np}\end{aligned}$$

Improvement?

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \leq -tnp] &\leq \inf_{\theta \leq 0} e^{tnp\theta} \prod_i \mathbb{E}[e^{\theta(X_i - \mathbb{E}[X_i])}] \\ &= \inf_{\theta \leq 0} e^{tnp\theta} e^{-np\theta} (1 + (e^\theta - 1)p)^n \\ &\leq \inf_{\theta \leq 0} e^{tnp\theta} e^{-np\theta} e^{np(e^\theta - 1)} \\ &= \exp\left(-np + \inf_{\theta \leq 0} (t - 1)np\theta + npe^\theta\right)\end{aligned}$$

- The minimizer is at $\theta = \log(1 - t)$:

$$\begin{aligned}\Pr[Z - \mathbb{E}[Z] \leq -tnp] &\leq \exp(-np - (1 - t)np \log(1 - t) + (1 - t)np) \\ &= \exp(-(1 - t)np \log(1 - t) - tnp) \\ &= \left(\frac{e^{-t}}{(1 - t)^{1-t}}\right)^{np}\end{aligned}$$

Chernoff Bound (Multiplicative Form)

Theorem (Chernoff bound).

Let X_1, \dots, X_n be iid Bernoulli random variables with $\mathbb{E}[X_i] = p$. Let $Z := \sum_{i=1}^n X_i$. Then

$$\Pr[Z \geq (1+t)np] \leq \left(\frac{e^t}{(1+t)^{1+t}} \right)^{np} \quad \forall t \geq 0$$

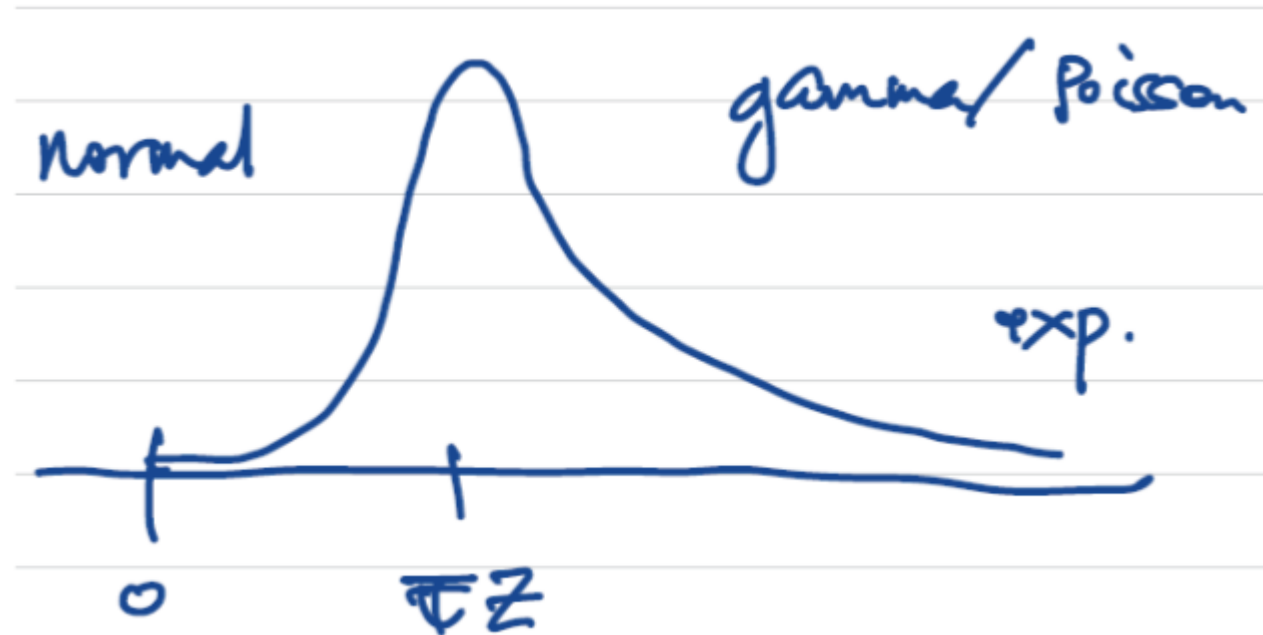
$$\Pr[Z \leq (1-t)np] \leq \left(\frac{e^{-t}}{(1-t)^{1-t}} \right)^{np} \quad \forall t \in [0,1]$$

Notice the asymmetric tails:

- For the right tail, when t is large, it behaves like t^{-t} (i.e., a Gamma distribution)
- For the left tail, when t is small, it behaves like $e^{-t^2/2}$ (i.e., a Gaussian distribution)

$$\log(\dots) = -t - (1-t) \log(1-t) = -t - (1-t)(-t - t^2/2 + \mathcal{O}(t^3)) = -t^2/2 + \mathcal{O}(t^3)$$

Chernoff Bound (Multiplicative Form)



(Source: Joel A. Tropp)

- $Z \sim \text{Bin}(n, p) \rightarrow \text{Poi}(np)$ for large n
- The probability of a Poisson random variable $X \sim \text{Poi}(np)$ at $k > np$ is

$$\frac{e^{-np}(np)^k}{k!} = e^{-\Theta(k \log k)} < e^{-\Omega(k^2)}$$

Chernoff Bound (Multiplicative Form)

Theorem (Chernoff bound, user-friendly version).

Let X_1, \dots, X_n be **independent Bernoulli** random variables with $\mathbb{E}[X_i] = p_i$. Let $Z := \sum_{i=1}^n X_i$ and $\mu := \mathbb{E}[Z]$. Then

$$\Pr[Z \geq (1 + t)\mu] \leq e^{-\frac{t^2}{2+t}\mu} \quad \forall t \geq 0$$

$$\Pr[Z \leq (1 - t)\mu] \leq e^{-t^2\mu/2} \quad \forall t \in [0,1]$$

Moreover,

$$\Pr[|Z - \mu| \geq t\mu] \leq 2e^{-t^2\mu/3} \quad \forall t \in [0,1]$$

Chernoff Bound (Multiplicative Form)

Theorem (Chernoff bound, user-friendly version).

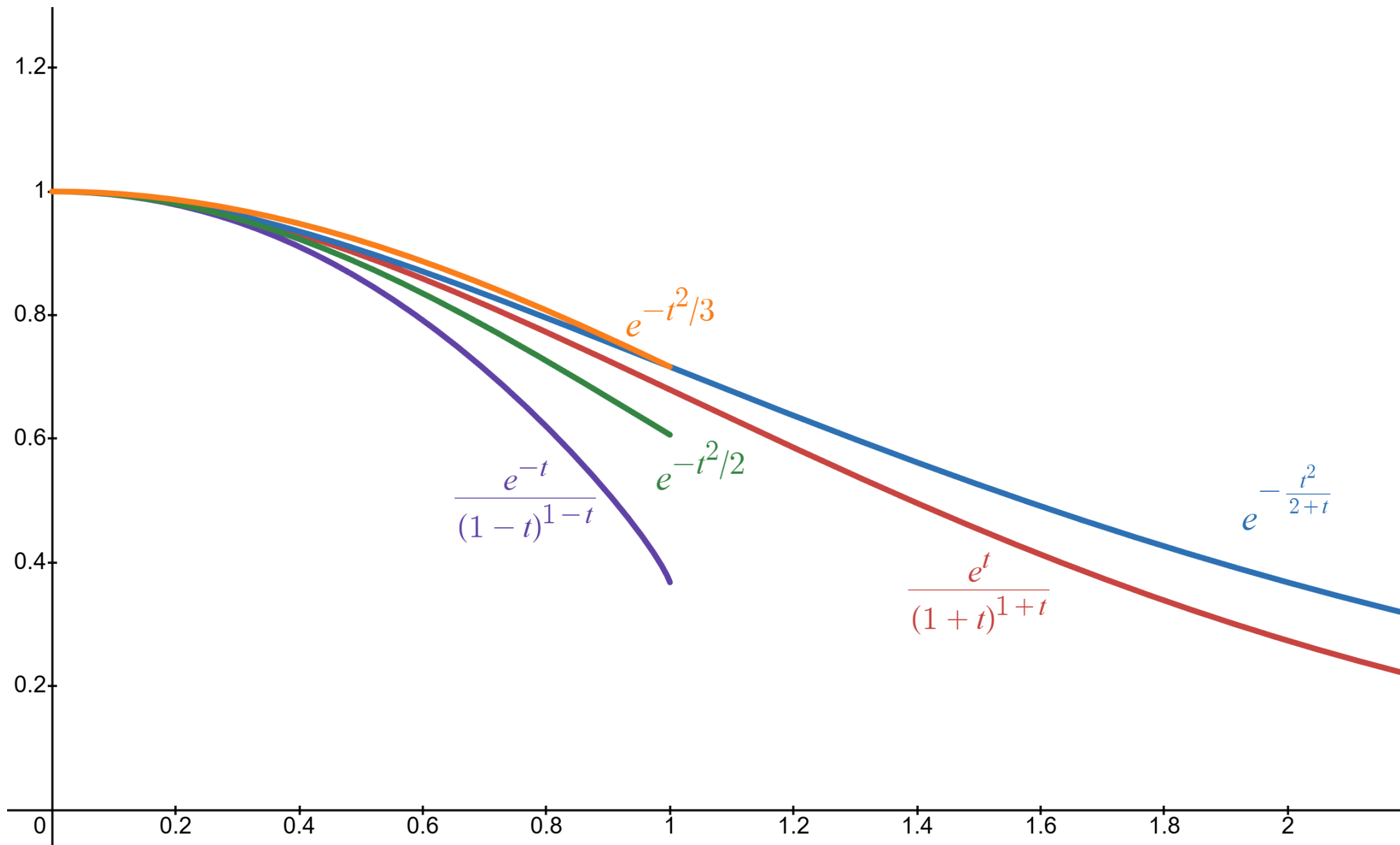
Let X_1, \dots, X_n be **independent Bernoulli** random variables with $\mathbb{E}[X_i] = p_i$. Let $Z := \sum_{i=1}^n X_i$ and $\mu := \mathbb{E}[Z]$. Then

$$\Pr[Z \geq (1 + t)\mu] \leq e^{-\frac{t^2}{2+t}\mu} \quad \forall t \geq 0$$

Proof.

- $\log\left(\frac{e^t}{(1+t)^{1+t}}\right) = t - (1+t)\log(1+t)$
- $\log(1+t) \geq \frac{2t}{2+t}$
- $t - (1+t)\log(1+t) \leq t - (1+t)\frac{2t}{2+t} = \frac{-t^2}{2+t}$





Chernoff Bound (Multiplicative Form)

Theorem (Chernoff bound, user-friendly version).

Let X_1, \dots, X_n be **independent Bernoulli** random variables with $\mathbb{E}[X_i] = p_i$. Let $Z := \sum_{i=1}^n X_i$ and $\mu := \mathbb{E}[Z]$. Then

$$\Pr[|Z - \mu| \geq t\mu] \leq 2e^{-t^2\mu/3} \quad \forall t \in [0,1]$$

➤ For our biased coin question,

$$\Pr[|Z - np| \geq (\epsilon/p)np] \leq 2 \exp\left(-\frac{\epsilon^2 n}{3p}\right) \leq \delta \quad \Rightarrow \quad n \geq \frac{3p}{\epsilon^2} \log\left(\frac{2}{\delta}\right)$$

Chernoff Bound (Multiplicative Form)

Theorem (Chernoff for positive, bounded random variables).

Let X_1, \dots, X_n be **independent** random variables such that $0 \leq X_i \leq b$ for all $i \in [n]$. Let $Z := \sum_{i=1}^n X_i$ and $\mu := \mathbb{E}[Z]$. Then

$$\Pr[Z \geq (1 + t)\mu] \leq \left(\frac{e^t}{(1 + t)^{1+t}} \right)^{\mu/b} \quad \forall t \geq 0$$

$$\Pr[Z \leq (1 - t)\mu] \leq \left(\frac{e^{-t}}{(1 - t)^{1-t}} \right)^{\mu/b} \quad \forall t \in [0, 1]$$

Proof of Chernoff for Positive Bounded RVs

- When X is a Bernoulli random variable,

$$\mathbb{E}[e^{\theta X}] = 1 + (e^{\theta} - 1)p$$

- If we only know $0 \leq X \leq b$, then by the convexity of $e^{\theta x}$ and Jensen inequality,

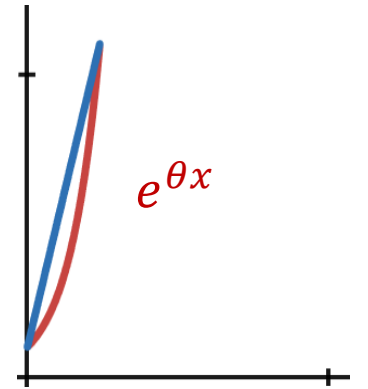
$$e^{\theta x} \leq 1 + (e^{\theta b} - 1)x/b \quad \forall x \in [0, b]$$

- Thus,

$$\mathbb{E}[e^{\theta X}] \leq 1 + \frac{e^{\theta b} - 1}{b} \mathbb{E}[X]$$

- Following the same proof, you'll get the term involving

$$\begin{aligned} \prod_i \left(1 + \frac{e^{\theta b} - 1}{b} \mathbb{E}[X_i] \right) &= \exp \left(\sum_i \log \left(1 + \frac{e^{\theta b} - 1}{b} \mathbb{E}[X_i] \right) \right) \leq \exp \left(\sum_i \frac{e^{\theta b} - 1}{b} \mathbb{E}[X_i] \right) \\ &= \exp \left(\frac{e^{\theta b} - 1}{b} \mu \right) \end{aligned}$$



Chernoff Bound (Additive Form)

Theorem (Chernoff-Hoeffding inequality).

Let X_1, \dots, X_n be iid Bernoulli random variables with $\mathbb{E}[X_i] = p$. Let $Z := \sum_{i=1}^n X_i$. Then

$$\Pr[Z \geq (p + \epsilon)n] \leq e^{-nD(p+\epsilon\|p)} \quad \forall \epsilon \in (0, 1 - p)$$

$$\Pr[Z \leq (p - \epsilon)n] \leq e^{-nD(p-\epsilon\|p)} \quad \forall \epsilon \in (0, p)$$

$D(p\|q)$ is the relative entropy defined as:

$$D(p\|q) = p \log \left(\frac{p}{q} \right) + (1 - p) \log \left(\frac{1 - p}{1 - q} \right)$$

Proof idea:

- Apply the Laplace transform and use the closed-form of MGF (same as our proofs)
- No more approximation and directly optimize it

Application: Discrepancy Theory

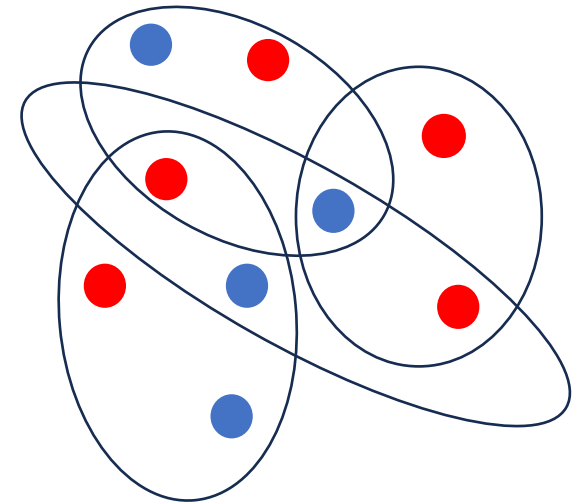
Given a hypergraph (or a set system), how to color each vertex in red or blue such that in each hyperedge (or set), the number of red vertices is roughly equal to the number of blue vertices?

- Universe $U = \{1, 2, \dots, n\}$
- Set system $\mathcal{F} = \{S_1, S_2, \dots, S_m\}$
- For a coloring $\chi: [n] \rightarrow \{-1, 1\}$, its discrepancy is defined as:

$$\text{Disc}(\chi) := \max_{i \in [m]} \left| \sum_{j \in S_i} \chi(j) \right|$$

- And the discrepancy of the set system is defined as

$$\text{Disc}(\mathcal{F}) := \min_{\chi} \text{Disc}(\chi)$$



$$\text{Disc}(\mathcal{F}) = 1$$

Application: Discrepancy Theory

Lemma. Let \mathcal{F} be a collection of m subsets of $[n]$. Then there is a coloring $\chi: [n] \rightarrow \{-1, 1\}$ such that $\text{Disc}(\chi) = \mathcal{O}(\sqrt{n \log m})$.

Proof.

- Let χ be a uniformly random coloring, i.e., we assign ± 1 to each element iid uniformly at random
- $\mathbb{E}[\chi(i)] = 0$ for any $i \in [n]$
- For each subset $S \in \mathcal{F}$, let $Z := \sum_{i \in S} \chi(i)$. By Hoeffding's inequality,

$$\Pr[|Z| \geq t\sqrt{v}] \leq 2e^{-t^2/2}, \quad v = \frac{1}{4} \sum_{i \in S} (1 - (-1))^2 = |S| \leq n$$

$$t := 2\sqrt{\log m} \quad \Rightarrow \quad \Pr[|Z| \geq 2\sqrt{n \log m}] \leq 2e^{-2 \log m} = 2/m^2$$

- By union bound over all m subsets, with probability $\geq 1 - m \cdot 2/m^2 > 0$, $\text{disc}(\chi) \leq 2\sqrt{n \log m}$



Application: Discrepancy Theory

Lemma. Let \mathcal{F} be a collection of m subsets of $[n]$. Then there is a coloring $\chi: [n] \rightarrow \{-1, 1\}$ such that $\text{Disc}(\chi) = \mathcal{O}(\sqrt{n \log m})$.

- This bound is not tight!

Theorem (Six standard deviations suffice, Spencer '85).

Let \mathcal{F} be a collection of n subsets of $[n]$. Then there is a coloring $\chi: [n] \rightarrow \{-1, 1\}$ such that $\text{Disc}(\chi) \leq 6\sqrt{n}$.

More generally, if \mathcal{F} is a collection of $m \geq n$ subsets of $[n]$. Then there is a coloring $\chi: [n] \rightarrow \{-1, 1\}$ such that $\text{Disc}(\chi) \leq \mathcal{O}(\sqrt{n \log(2m/n)})$.

Bernstein Inequality

Theorem (Bernstein).

Let X_1, \dots, X_n be independent random variables such that $|X_i - \mathbb{E}[X_i]| \leq b$ for all $i \in [n]$.
Let $Z := \sum_{i=1}^n X_i$. Then

$$\Pr[|Z - \mathbb{E}[Z]| \geq t] \leq 2 \exp\left(-\frac{t^2/2}{\text{Var}[Z] + bt/3}\right) \quad \forall t > 0$$

- For large t , the tail bound looks like $\exp(-2t/(3b))$, an exponential tail
- For small or medium t , the tail bound looks like $\exp(-t^2/(2\text{Var}[Z]))$, a Gaussian tail with the true variance (instead of the variance proxy v in Hoeffding's inequality)